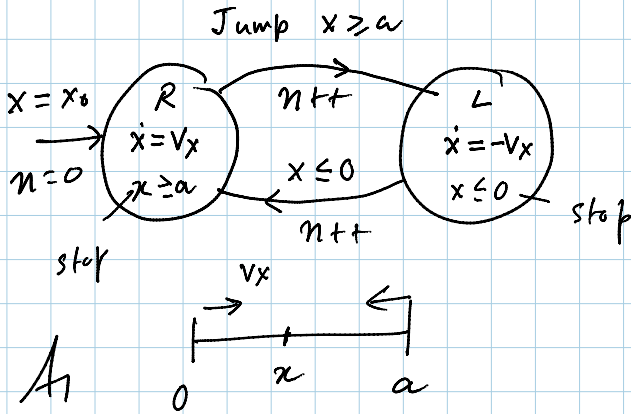
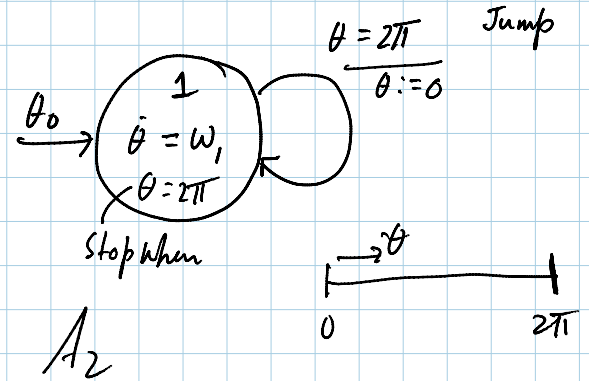


# HW1 Problem 2

## 1D billiard ball



## 1D - Satellite



$A_1$  &  $A_2$  are bisimilar

$A_1$  implements  $A_2$  and vice versa

There exists a Forward Sim Rel from  $A_1$  to  $A_2$

$R_1$ : Candidate FSR  $R_1 \subseteq \mathcal{Q}_1 \times \mathcal{Q}_2$

$$R_1 \triangleq \left\{ (s, u) \mid s \in \mathcal{Q}_1, u \in \mathcal{Q}_2, s.x = \frac{a}{2\pi} u.\theta \right\}$$

$\text{type}(x) = [0, a]$   
 $\text{type}(\text{loc}) = \{R, L\}$

$\mathcal{Q}_1 \stackrel{P}{=} \text{val}(\{x, \text{loc}\})$

$\mathcal{Q}_2 \stackrel{A}{=} \text{val}(\{\theta, \text{loc}_2\})$

Rough calc

$$\frac{x}{a} = \frac{\theta}{2\pi}$$

$$x = \frac{a}{2\pi} \theta$$

$$v_x = \frac{a}{2\pi} \omega$$

$$x_0 = \frac{a}{2\pi} \theta_0$$

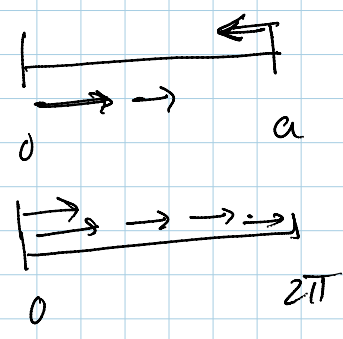
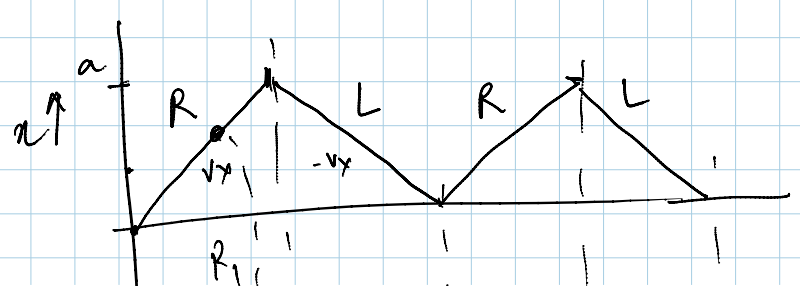
Assumptions

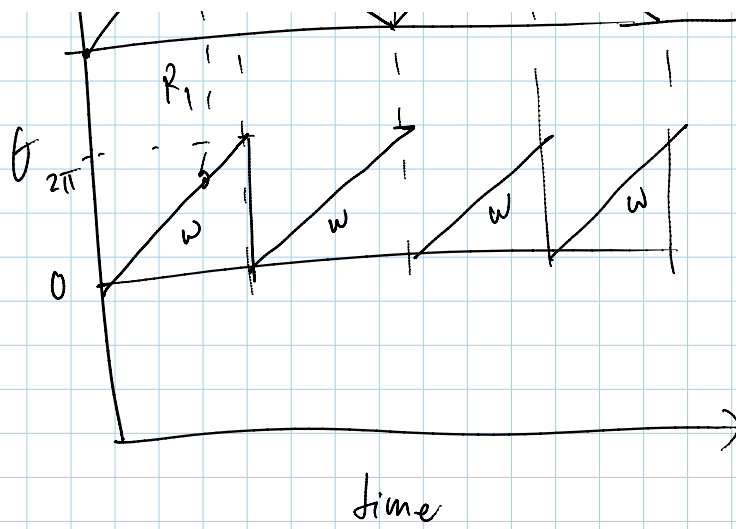
(1)

(2)

Claim:  $R_1$  is a FSR from  $A_1$  to  $A_2$ .

Proof:



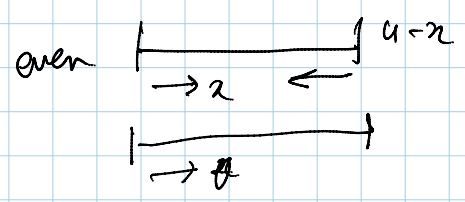


$$Q_1 = \text{val}(\{x, x, \text{loc}\})$$

$$Q_2 = \text{val}(\{\theta, \text{loc}\})$$

$$R_2 = \left\{ (s, u) \mid \begin{array}{l} s, x = \text{even} \Rightarrow \frac{a}{2\pi} \theta = x \\ s, x = \text{odd} \Rightarrow \frac{a}{2\pi} \theta = a - x \end{array} \right\} \quad \text{--- (3)}$$

$$\text{--- (4)}$$



Proof  
Base

inductive Step

(a)  $s \xrightarrow{\gamma_1} s'$  and  $s R_2 u$   
we show  $\exists \gamma_2, u'$  s.t.  $u \xrightarrow{\gamma_2} u' \wedge s' R u'$

$\gamma_2$  uniquely defined by  $u$  and  $\gamma_2.\text{time}$   
 $\gamma_2.\text{time} \triangleq \gamma_1.\text{time}$ , say  $= t$

$$\left. \begin{array}{l} s'.x = s.x + v_x \cdot t \\ u'.\theta = u.\theta + w_c \cdot t \end{array} \right\} \text{by def of } A_1, A_2$$

Case (i)  $s, x = \text{even}$

$$u'.\theta = u.\theta + w_c \cdot t$$

$$= \frac{2\pi}{a} s.x + w_c \cdot t \quad \left[ \text{From } R_2 \text{ (3)} \right]$$

Assumption

$$\begin{aligned}
 &= \frac{2\pi}{a} s \cdot x + w_1 t \\
 &= \frac{2\pi}{a} s \cdot x + \frac{2\pi}{a} v_x \cdot t \quad [\text{From Assumption (1)}] \\
 &= \frac{2\pi}{a} [s \cdot x + v_x \cdot t] \\
 &= \frac{2\pi}{a} [s' \cdot x]
 \end{aligned}$$

$s' \cdot x = \text{even}$  so  $w'$  and  $s'$  satisfies  $R_2(3)$

Case (ii)  $s \cdot x = \text{odd}$

Next we show that  $\gamma_2$  is indeed allowed from  $u$ .

$\forall t' \in \gamma_2 \cdot \text{dom}$   $\gamma_2(t')$  does not satisfy the

stopping cond  $0 \leq \theta < 2\pi$

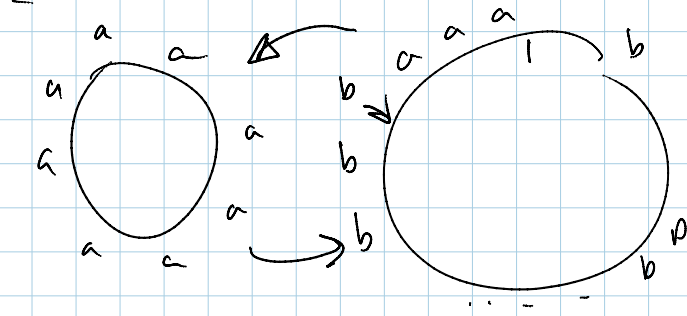
$$\gamma_2(t') \cdot \theta = u \cdot \theta + t' \cdot w_1$$

$$= \frac{2\pi}{a} (s \cdot x + t' \cdot v_x) \quad [\text{using } R_2]$$

$$< \frac{2\pi}{a} \cdot a = 2\pi$$

(b)  $s \xrightarrow{\text{bounce}} s'$  and  $s R u$  band

we show  $\exists u'$  s.t.  $u \rightarrow u' \text{ s' R } u'$



## Stability

Hybrid automaton  $A$

Recall, an execution of  $A$  is alternating sequence

$$\alpha = \tau_0 a_1 \tau_1 a_2 \dots$$

$$\tau_i : [0, t] \rightarrow \text{val}(x)$$

$$a_i \in A$$

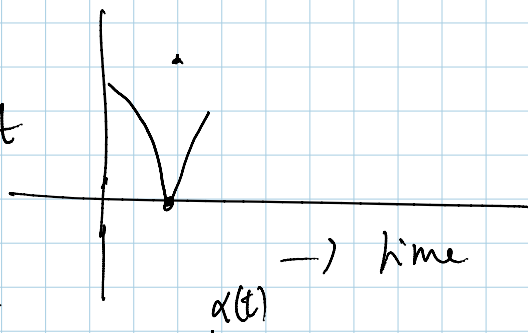
$$\tau_i \cdot \text{state} \rightarrow \tau_{i+1} \cdot \text{state}$$

$$\alpha(t) \quad t \in [0, \alpha.\text{time}]$$

$$\alpha.\text{time} = \sum \tau_i.\text{time}$$

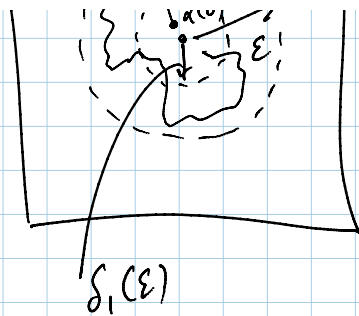
$$\alpha(t) \stackrel{\Delta}{=} \alpha' \cdot \text{state}$$

where  $\alpha'$  is the longest prefix of  $\alpha$  with  $\alpha' \cdot \text{time} \leq t$



Stability is a property of





Remarks If  $\delta_1(\epsilon) \rightarrow 0$  then  $\epsilon$ -Ball invariants for free